*Original Article*

# Hybrid Hardware in Loop Architecture for Autonomous Vehicles

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*Abstract* - *The automotive world is going through a paradigm shift. With sensor technologies like 4K cameras, LiDARs, Radars, etc., becoming more affordable and advances in Deep Learning algorithms development, research is trending towards developing Autonomous Vehicles and bringing them mainstream. However, a challenge remains in the commercialization of technologies. An automotive vehicle undergoes a rigorous verification and validation model. Perfected over decades, standards and guidelines are introduced to minimize failures. However, there exists a gap in testing when working with new sensors and algorithms, and our paper focuses on a hybrid setup for hardware in loop testing, which enables software to iterate over their development faster without having to wait to deploy on actual cars for testing.*

*Keywords - Autonomous Vehicles, Linux, Operating System, Emulation, QEMU.*

# **1. Introduction**

Embedded Machine Learning, computing on the edge is crucial for fostering innovation in a multitude of intelligent devices. The evolution of ARM processors has created an ecosystem of smart devices where significant computing exists to run Machine Learning Models. The next innovation is touted to be in the automotive sector, where machine learning and hardware combine to improve the transportation sector. The edge computer means that the decision-making power is now with an ECU [1] (Engine Control Unit). A car can have multiple ECUs, with each having its own functionality, from controlling the wipers to controlling the steering accelerator based on the user input. In autonomous vehicles, a software agent shall drive these actuators. This makes the testing of the software, along with the underlying hardware, a challenge. Hardware-in-the-loop (HIL) testing involves connecting a real hardware system (such as an ECU) to a virtual environment that simulates real-world conditions.

This allows engineers to test the control systems' behavior under various conditions without needing the entire vehicle. HIL testing allows identify potential issues early in the development process, thus reducing the need for expensive and time-consuming physical prototypes. Figure 1 shows a V-Model, also called the verification and validation model. International industry standards prescribe the model for the development of safety-critical systems like ISO26262. While the left side of the V model focuses on requirements at both macro and micro levels, the right side focuses on software integration and testing. In the initial stages, the software is tested in a purely simulated environment to allow for unit and module testing, while the later stages test the overall software on the target hardware platform with a full load of sensors.

Often, the testing infrastructure between the initial stages in the simulated environment vs the actual target hardware brings in a huge variation. For instance, a software driver is implemented to capture 3d point cloud data from a LiDAR. The software driver in order to test, will fake add function stubs in code to mimic the sequence of operations to be performed on the LiDAR hardware only to see hardwarerelated issues when testing the software on the target hardware.

Our paper aims to bridge the gap in the testing process with a Hybrid Hardware Loop Testing process, which uses QEMU [2] virtualization and creating custom device models in QEMU to enable hardware systems to be patched through to the virtualized environment.

The rest of the paper is organized as follows. Introduction to hardware emulators in section II. Section III describes QEMU, Section IV describes the Process of Adding User Defined Devices in QEMU. Section V shall have the Device Driver Development in Linux and we shall conclude in section VI. Fig. 1 v- model of development





cd gemu mkdir build cd build ../qemu/configure --target-list=aarch64-softmmu make

### **Fig. 2 Qemu build instruction**

qemu-system-aarch64 -M virt -cpu cortex-a57 \ -smp 4 -m 1G -kernel Image  $\setminus$ -append "root=/dev/vda rw console=ttyAMA0" \ -semihosting -nographic -drive file=rootfs.ext4

**Fig. 3 Invoking QEMU**

## **2. Hardware Emulators**

An embedded system product development depends on the availability of the hardware. In the absence of hardware and for the software development to progress, developers often use hardware emulators. For instance Android developers use the Android Emulator [3]for developing their application. In this context, emulation can be defined as using a host CPU's resources to mimic the functionality of another CPU architecture. Thus a software developed for say ARM CPU architecture can be run on a host system which has an x86 architecture. This is the typical use-case of an Android emulator for android application development. Some widely used emulation products are QEMU, BOCHS [4] RENODE [5] In automotive ECU HIL testing commercial setups like dSPACE [6] are often used for validation. Our paper focuses on QEMU and its application in the Hardware in Loop Testing. QEMU, an open-source emulator which employs dynamic binary translation to achieve high performance when emulating guest systems on different host architectures. The modular nature of QEMU allows it to emulate a wide range of hardware components, including CPUs, network interfaces, and storage controllers, making it highly adaptable for various use cases. QEMU's KVM (Kernel-based Virtual Machine) integration allows for near-native execution speeds by leveraging hardware virtualization extensions if supported by the host system. While QEMU supports multitude of standard hardware interfaces, the Autonomous Vehicle domain brings in new hardware which are either not modeled in QEMU or not yet available in public domain to be accessible to everyone.

## **3. QEMU**

We will need the QEMU source code to add custom hardware. You can download the source from the official QEMU repository https://gitlab.com/qemu-project/qemu.git. To build QEMU source code: We also require a root file system and a Linux Kernel Image to use on our custom QEMU emulator. You can use the debootstrap [7] tool to generate a Debian [8] base image and the Linux Kernel [9] source code from the link below. https://git.kernel.org/pub/scm/linux/kernel/git/stable/linux.git

Another popular option is to use buildroot [10] [\(https://buildroot.org/\)](https://buildroot.org/) for generating a root file system along with the kernel image. Since we are generating an image for the QEMU Virtual Platform, we shall use the config *qemu\_aarch64\_virt\_defconfig.* Once both qemu and buildroot finish the *making* process, Qemu can be invoked using the command below, shown in Figure 3. QEMU shall execute with the kernel image and the roots parameter to start a serial console emulating an ARM64 [11] virtual system, as shown in Figure 4.

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# **4. User Defined Custom Hardware Model**

With the standard QEMU invocation completed, we shall now focus on creating a custom hardware model. To re-iterate, this user-defined model being developed can be used to replicate a hardware interface, say a sensor register layout, add function hooks into a world view simulator CARLA [12] or add software to create a pass-through of the camera data, which is connected to the host via a special interface like GMSL [13]. Our target system is an Embedded ARM64 SOC, as shown in Figure 5, with a user-defined co-processor added to the system bus. To show QEMU's scalable and modular architecture, we can even send an interrupt event from this coprocessor to the main Application Processor. In the QEMU source code directory, we shall modify the virtual platform defined under *hw/arm/virt.c* and its associated header file*.* In the *include/hw/arm/virt.h* file we add a device type *VIRT\_CUSTOM\_COPROC*. ARM/ARM64 architecture has a flat memory model.





The data structure *MemMapEntry* holds the memory map of the entire SOC. We shall add an entry in a vacant space at address 0x0b000000 with a length of 0x200. Similarly, we add a virtual IRQ from our user-defined co-processor to the Application Processor. Next, we shall create the software model of our user-defined model, which is slotted at base address 0x0b000000 and has an associated interrupt. Under include/hw/arm and under hw/arm, create a new sub-folder, which will contain the header file and the source file of your model.Our user-defined model has a device id of *0x42024*. The id offset can be read at base\_address + 0x0 (kPeripheralIDOffset). The simple model shown here shall take in two 32-bit input numbers in the input registers from the Application Processor and shall return the output of it in the output register. The register at offset 0x14 (kStartProcessing) shall be the trigger from the software driver to begin the compute. Additionally, the device driver can check if the user modeled coprocessor is busy with the register at 0x10 (kProcIdle) as well as reset the co-processor with a reset register at 0x30 (kResetCoProc). If the Device Driver

configures the interrupt (register offset 0x8) to be enabled, the user model shall send a completion interrupt back to AP once the execution is completed.

static const MemMapEntry base memmap[] = {	
	/* Space up to 0x8000000 is reserved for a boot ROM */
$[VIRT FLASH] =$	$0.008000000$ }.
$[VIRT CPUPERIPHS] =$	$\{ 0 \times 08000000, 0 \times 00020000 \}.$
$[VIRT MMIO] =$ $[VIRT$ CUSTOM COPROC] =	$\{ 0x0a000000, 0x00000200 \}.$ $\{ 0 \times 00000000, 0 \times 000000200 \}.$
$[VIRT MEM] =$	/* Actual RAM size depends on initial RAM and device memory settings */ { GiB, LEGACY RAMLIMIT BYTES },

**Fig. 6 Virtual ARM Platform Memory Map**

```
static const int a15irqmap[] = {
    [VIRT_UART] = 1,[VIRT_RTC] = 2,
    [VIRT_SMMU] = 74,/* ...to 74 + NUM_SMMU_IRQS - 1 */
    [VIRT_PLATFORM_BUS] = 112, /* ...to 112 + PLATFORM_BUS_NUM_IRQS -1 */
    [VIRT_CUSTOM_COPROC] = 112 + PLATFORM_BUS_NUM_IRQS,
\mathbf{E}
```


} CoProcessorState;

#### **Fig. 8 User defined model attributes**

Our new model registration with QEMU uses the macro type\_init.

```
static const TypeInfo _coProcessorInfo = {
     .name = TYPE CUSTOM COPROC.
     .parent = TYPE_SYS_BUS_DEVICE,
     .instance_size = sizeof(CoProcessorState),
     .instance_init = coProcessorInit,
     .instance_finalize = coProcessorCleanUp,
\mathcal{H}static void registerCoProcessor(void)
-{
    type_register_static(&_coProcessorInfo);
\mathcal{L}
```
type\_init(registerCoProcessor)

<sup>1</sup>

#### **Fig. 9 Code to register the user-defined model**

```
static const MemoryRegionOps CoProcOps = {
    .read = CoProcRead,
     write = CoProcWrite.endianness = DEVICE_NATIVE_ENDIAN,
1:
static void CoProcessorWorkCompletion(void *opaque)
-1
    CoProcessorState *s = (CoProcessorState *)opaque;
    s\rightarrow out = s\rightarrow input + s\rightarrow input;s->busy = false;
    s->irqStatus = true;
    CoProcSetIRO(s);
\mathbf{v}static void coProcessorInit(Object *obj)
    CoProcessorState *state = CoProcessor(obj);
    SysBusDevice *sbd = SYS_BUS_DEVICE(obj);
    memory_region_init_io(&state->iomem, OBJECT(state), &CoProcOps, state,
                           TYPE CUSTOM COPROC, 0x200);
    sysbus_init_mmio(sbd, &state->iomem);
    sysbus_init_irq(sbd, &state->irq);
    // initialize internal state
    state->id = kPeripheralID;
```

```
state->timer = timer_new_ns(QEMU_CLOCK_VIRTUAL, CoProcessorWorkCompletion, state);
```
#### **Fig. 10 Init**

```
static uint64_t CoProcRead(void *opaque, hwaddr offset, unsigned size)
    CoProcessorState *s = (CoProcessorState *)opaque:
    uint32 t ret = 0xDEADEEF:switch (offset)
    case kPeripheralIDOffset:
        ret = s - sid;break:
    ....<br>case kCoProcIdle:
        ret = !s->busy;<br>break;
```
case kOutputReg:  $ret = s$ ->out:  $break:$ 

return ret: Ï

#### **Fig. 11 MMIO read function**

static void CoProcWrite(void \*opaque, hwaddr offset, uint64\_t value, unsigned size)

```
ł
    CoProcessorState *s = (CoProcessorState *)opaque;
    switch (offset)
    case kCompletionIRQEnOffset:
        s->en = value & 0x1;
        break;
    case kInpReg1:
       s->inp1 = value;
    break:
    case kInpReg2:
       s->inp2 = value;
    break;
    break:
    case kStartProcessing:
       s->busy = value ? true : false;
        if (s->busy) {
            s->out = 0;
            setUpTimer(s):
        break;
   - 1
\mathbf{I}
```
#### **Fig. 12 MMIO write function**

static void CoProcessorWorkCompletion(void \*opaque)

```
CoProcessorState *s = (CoProcessorState *)ovalue:s->out = s->inp1 + s->inp2;
    s->busy = false;
    s->irqStatus = true:
    CoProcSetIRQ(s);
\mathbf{I}
```
#### **Fig. 13 QEMU timer callback**

The code snippets in Figures 9 to 14 show how the coprocessor is registered with QEMU, its init sequence where the memory device memory region is intialized with the length 0x200. The read/write APIs get invoked when a driver makes a hardware register access in the co-processor assigned memory section. Once the driver sets a non-zero value in the kStartProcessing register, we trigger the QEMU internal timer. On completion of the time period, we receive a callback where we output the result of the execution as well as send an interrupt to the Application Processor. After adding the code and adding the source files to the QEMU Kconfig/Meson build system, we are ready to rebuild the QEMU binary.

```
static void create custom coproc(const VirtMachineState *vms)
     hwaddr base = vms->memmap[VIRT_CUSTOM_COPROC].base;<br>hwaddr size = vms->memmap[VIRT_CUSTOM_COPROC].size;<br>int irq = vms->irqmap[VIRT_CUSTOM_COPROC];
     char *nodename;
    sysbus_create_simple(TYPE_CUSTOM_COPROC, base, qdev_get_gpio_in(vms->gic, irq));
     MachineState *ms = MACHINE(vms);
     nodename = g_strdup_printf("/coproc@%" PRIx64, base);<br>qemu fdt add subnode(ms->fdt, nodename);
     qenu_fdt_setprop_string(ms->fdt, nodename, "compatible", "coproc");
     qemu_fdt_setprop_sized_cells(ms->fdt, nodename, "reg",
     quantity and the proposition of the proposition of the proposition of the set of the proposition of the set of the CIC_FDT_IRQ_TYPE_SPI, irq,
                                     GIC_FDT_IRQ_FLAGS_LEVEL_HI);
     g_free(nodename);
\mathbf{A}Fig. 14 Instantiating the CoProcessor in virt.c
```
# **5. Device Driver and Validation**

With the new compiled qemu binary, if we start the emulator with the ARM virtual machine in the monitor mode and run the command info mtree, you shall see the userdefined coprocessor in the system address map.

Next, if we read the memory using devmem [14] [at the base-address of our new device, we should read the hardcoded device id, which confirms our user-model read operation is working as expected.

```
build ./qemu-system-aarch64 -machine virt -monitor stdio
QEMU 9.0.50 monitor - type 'help' for more information
                   (qemu) info mtree
             address-space: cpu-memory-0
                address-space: memory
   0000000000000000-ffffffffffffffff (prio 0, i/o): system
     0000000009010000-0000000009010fff (prio 0, i/o): 
                         pl031
                           ....
                            ....
   0000000000000000-0000000000000001ff (prio 0, i/o):
                 CoProcessor-Custom
```
**Fig. 15 Address map with user model**

# devmem 0xb000000 0x00042024 #

We shall write a tiny linux device driver which will match against the compatible string "coproc" as defined earlier in our hw/arm/virt.c file. We read the interrupt property in the device tree and register an interrupt handler for it. Recall that our coprocessor fires an interrupt when the device driver initiates the processing of the inputs by writing to the register at offset 0x14, which in turn triggers an internal QEMU Timer. The timer call back is used to send an interrupt to the AP.

```
static int coproc_remove(struct platform_device *pdev)
{
   return 0;
}
static const struct of device id coproc of match[ = {
  \{.compatible = "coproc", \},\ { }
};
MODULE DEVICE TABLE(of, coproc of match);
static struct platform driver coproc driver = {
  .probe = coproc probe,
  .remove = coproc_remove,
  .driver = \{ .name = "coproc-driver",
    .of match table = of_match_ptr(coproc_of_match),
   },
};
module_platform_driver(coproc_driver);
```
**Fig. 16 Registering the linux device driver**

```
Fig. 17 The probe function
static int coproc_probe(struct platform_device *pdev)
{
  struct device *dev = \& pdev->dev;
 ...
   cd->base = devm_ioremap(dev, res->start, 
resource size(res));
  cd-\veeirq = irq_of_parse_and_map(pdev-\geqdev.of_node,
(0):
  if (cd-\gevirq == 0) {
   }
   else {
     ret = request\_irq(cd \rightarrow virq,(irq_handler_t)coproc_irq_handler,
                IRQF_TRIGGER_RISING, 
"COPROC_IRQ_HANDLER", cd);
 }
   // Enables Interrupt
  writel(1, cd->base + 0x8);
}
```


**Fig. 18 The Co-Processor interrupt handler**

We compile the device driver into the Linux Kernel Image and re-invoke QEMU with the new kernel image. On running cat */proc/interrupts,* we should see our IRQ handler registered.



Welcome to Buildroot	
buildroot login: root	
$#$ sh test.sh	
CoProc ID:	
0x00042024	
CoProc Will Add 10 20	
Trigger Start Processing	
Sending IRO to Main Processor	
-CoProc Irg Rx	
CoProc Irg Cleared When Status Reg is Read	
Check Completion Status Reg	
0x00000000	
Outout Rea	
0x0000001E	
ᡰ#᠋	

**Fig. 20 Running the shell script**

## **References**

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We shall use a shell script with devmem to write two input numbers and observe the output. An interrupt is generated on completion of the work and we see the print statements in the serial console.

## **6. Conclusion**

In this paper, we present a new hybrid approach to performing Hardware Loop Testing for automotive applications. QEMU being a versatile emulator, we add a userdefined model using the 'C' programming language.

With the HIL setup as part of the test infrastructure, developers can focus on higher-level Computer Vision and Deep Learning problems with real input data such as camera images, radar data or the 3d point cloud data from a LiDAR.

For future work, we intend to provide create a sensor framework layer which will allow the user-defined model to interface with automotive sensors seamlessly.